

Quantum Sets

Gelfand Duality

All vector spaces are complex. A normed vector space is a vector space V with a norm function

$$\|\cdot\|: V \rightarrow \mathbb{R}^+$$

A norm induces a metric on V , and V is a Banach space if it is complete in this metric.

The category Ban is the category whose objects are Banach spaces and whose arrows are continuous linear maps.

Thm: Ban is monoidal via the projective tensor product. In fact, it's symmetric monoidal closed.

Defn: A Banach algebra is a Banach space A with an associative multiplication

$$m: A \otimes A \rightarrow A$$

Warning: Some sources require a unit. Some don't.
We'll say unital if we want a unit.

A C^* -algebra B is a Banach algebra with
an operation $(\cdot)^*: B \rightarrow B$ s.t.

$$x^{**} = x$$

$$(x+y)^* = x^* + y^*$$

$$(xy)^* = y^* x^*$$

$$(\lambda x)^* = \overline{\lambda} x^* \quad \lambda \in \mathbb{C}$$

$$\|x^* x\| = \|x\| \cdot \|x^*\|$$

Ex: 1) Let X be a compact Hausdorff space.

$$C(X) = \{ f: X \rightarrow \mathbb{C} \mid \text{continuous} \}$$

$$\|f\| = \sup \{ \|f(x)\| \mid x \in X \} \quad f^*(x) = \overline{f(x)}$$

is a unital commutative C^* -algebra.

2) Let X be locally compact Hausdorff

$$C_0(X) = \{ f: X \rightarrow \mathbb{C} \mid \text{continuous \& "vanishes at } \infty \}$$

3) If H is a Hilbert space,

$$B(H) = \{ T: H \rightarrow H \mid T \text{ is bounded \& linear} \}$$

31

$B(H)$ is a noncommutative unital C^* -Algebra

4) $K(H) = \{ f: H \rightarrow H \mid f \text{ is linear and compact} \}$

is a non-unital C^* -algebra

Let Cpt be the category of compact Hausdorff topological spaces.

Let $C^*-Alg_{c,u}$ be the category of commutative, unital C^* -algebras with unit, (and obvious notion of morphism)

$$C: Cpt^{op} \rightarrow C^*-Alg_{c,u}$$

There's a functor

$$sp: C^*-Alg_{c,u}^{op} \rightarrow Cpt$$

and these two functors form an equivalence

$sp(A) = \{ \chi: A \rightarrow \mathbb{C} \mid \chi \text{ is an algebra homomorphism} \}$. It can be given a topology that makes this work

Noncommutative Topology/Geometry,

Since commutative C^* -algebras can be thought of as functions on a space, then a noncommutative C^* -algebra can be thought of as functions on a noncommutative space. Can we redo geometry at this abstract level?

The theory of quantum sets is only finite for the moment.

Finite Gelfand duality

The category of finite-dimensional commutative C^* -algebras is equivalent to the opposite of the category of finite sets and functions.

We'd like a more abstract characterization which didn't rely on the existence of norms.

(5)

Compact Closed Categories (Max Kelly)

A symmetric monoidal ~~closed~~ category is compact closed if it comes equipped with an operation on objects

$$C \mapsto C^*$$

and for all objects, maps

$$I \xrightarrow{\varepsilon} C \otimes C^* \quad C^* \otimes C \xrightarrow{\eta} I$$

satisfying 2 equations. Here's one of them, ignoring associativity

$$\begin{array}{ccc}
 C & \xrightarrow{\sim} & I \otimes C \xrightarrow{\varepsilon \otimes \text{id}} C \otimes C^* \otimes C \\
 & \searrow \text{id} & \downarrow \text{com} \\
 & & C \otimes I \\
 & & \downarrow \sim \\
 & & C
 \end{array}$$

The other equation begins with C^* .

You can prove many things:

1) $()^*$ can be extended to a functor

$$()^* : \mathcal{C}^{op} \rightarrow \mathcal{C}$$

$$2) ()^{**} \simeq \text{id} : \mathcal{C} \rightarrow \mathcal{C}$$

3) A compact closed category is closed
with $A \Rightarrow B \triangleq A^* \otimes B$

$$4) (A \otimes B)^* \simeq A^* \otimes B^*$$

5) Endomorphisms in a compact closed
category have a trace

$$\underline{f : A \rightarrow A}$$

$$\text{Tr}(f) : I \rightarrow I$$

satisfying $\text{Tr}(fg) = \text{Tr}(gf)$

This leads to the theory of traced monoidal
categories

We can also talk about string diagrams
in compact closed category

$$\cap_{\alpha}: I \rightarrow X \otimes X^*$$

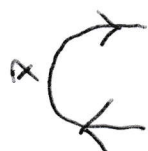
(7)

$$U_{\alpha}: X^* \otimes X \rightarrow I$$

We'll use the convention that on a X arrow, pointing left means X , but pointing right means X^* , and we'll draw our string diagrams left to right.

$$I \rightarrow X \otimes X^*$$

$$X^* \otimes X \rightarrow I$$



our first equation becomes:



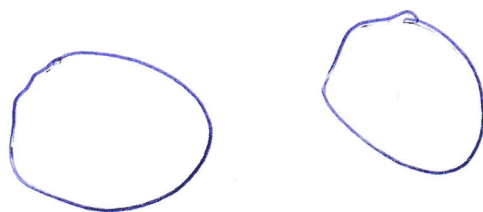
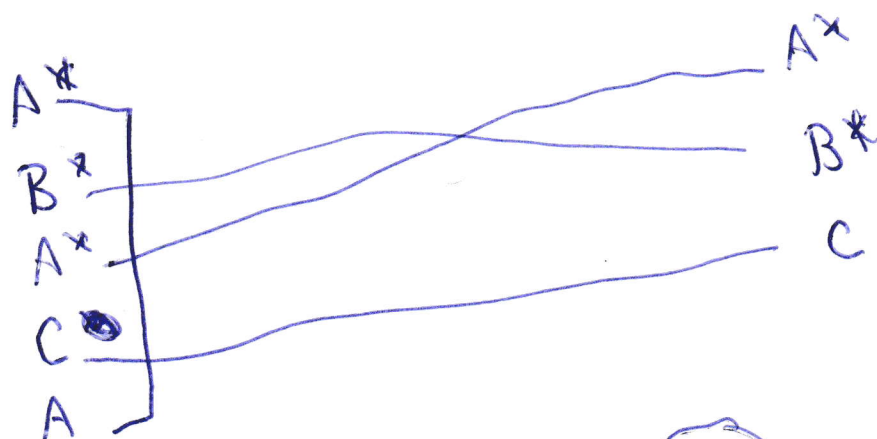
Most equations in monoidal categories can be expressed in terms of

Loops can be created in these string diagrams

$$I \rightarrow A \otimes A^* \quad A \otimes A^* \rightarrow I$$



So in the free compact closed category,
a typical morphism will be a string diagram with some loops



Kelly & Laplaza

this sort of diagram

②

COHERENCE THEOREMS

We're going to follow

A compositional approach to quantum foundations
by Musto, Reutter, Verdun

They're working with Hilbert spaces,
not just vector spaces. So there's additional
structure

$$(\)^{\dagger} : \text{Hilb}^{\text{op}} \rightarrow \text{Hilb}$$

which is an involution, and the identity
on objects

$$\begin{array}{c} f : H \rightarrow K \\ \hline f^{\dagger} : K \rightarrow H \end{array}$$

and

$$\begin{array}{c} f \otimes g : H_1 \otimes H_2 \rightarrow K_1 \otimes K_2 \\ \hline \end{array}$$

$f^{\dagger}$

Following the paper, we'll draw our diagrams, going upward

(9)

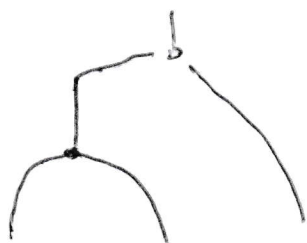


$$m: H \otimes H \rightarrow H$$



$$u: C \rightarrow H$$

For an associative, unital algebra these satisfy



=



=

|

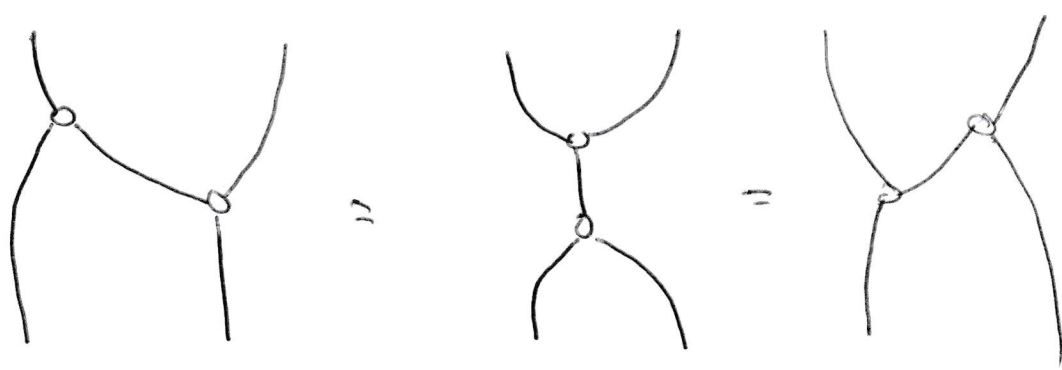
=



In a dagger category, one can automatically turn an algebra into a coalgebra. This is called the adjoint coalgebra structure

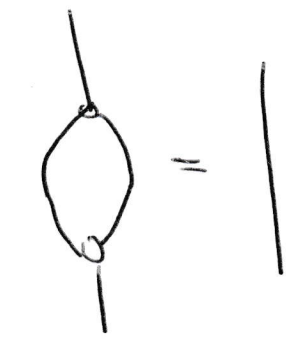
Def'n: A dagger Frobenius algebra

is an algebra in a monoidal dagger category such that

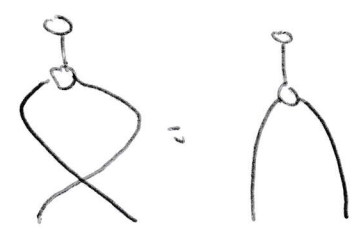


This is the Frobenius equation. In general, in a monoidal category, not necessarily dagger, a Frobenius algebra is an algebra & coalgebra satisfying this equation.

There are several types of Frobenius algebras



special



symmetric



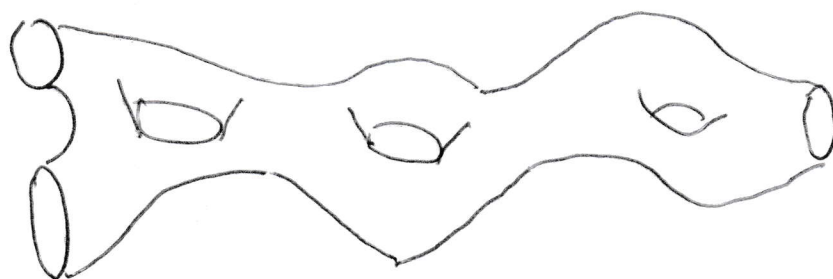
commutative

2D TQFT

⑪

See Joachim Kock - Frobenius Algebras
& 2D TQFT

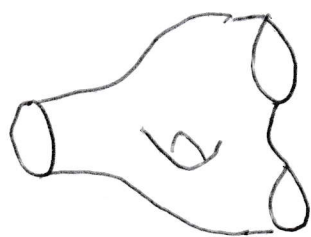
Build a category (roughly speaking)



$2 \rightarrow 1$

Objects are a finite disjoint union of circles,
arrows are cobordisms as above.

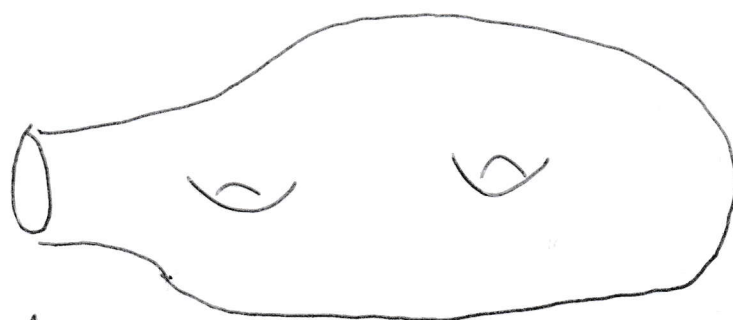
Composition is to glue along boundaries.



$1 \rightarrow 2$



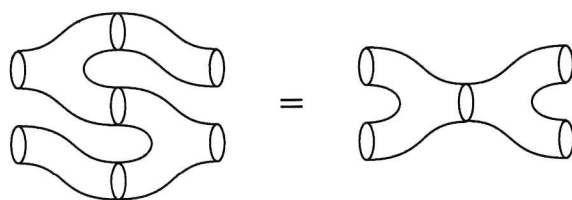
$2 \rightarrow 0$



1

$\rightarrow 0$

This is a
category
 2Cob



FRONTISPIECE *The topological expression of the main axiom of a Frobenius algebra. This figure illustrates the content of the main theorem, and captures the whole spirit of the book.*

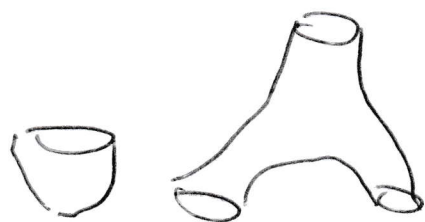
This is a monoidal category with disjoint union as $\otimes$

(12)

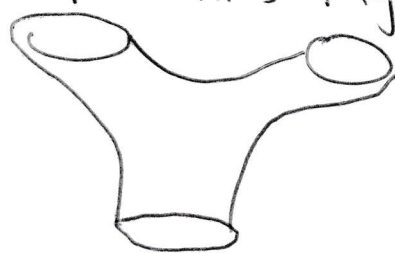
Def'n: A 2D TQFT is a symmetric monoidal functor from

$$F: 2\text{Cob} \rightarrow \underbrace{\text{Vec}_{\mathbb{C}}}_{\text{complex fd vector spaces}}$$

Note: In 2Cob , the object 1 (1-circle) has the structure of a Frobenius Algebra (commutative)



Algebra



Coalgebra



Thm: The category 2Cob is the symmetric monoidal category freely generated by the above commutative Frobenius algebra

Thm: There is an equivalence of categories

$$\begin{array}{ccc}
 2\text{TQFTs} & \simeq & \text{cFA} \\
 \nearrow & & \nwarrow \\
 \text{2d TQFTs} & & \text{commutative FAs} \\
 \text{and appropriate} & & \\
 \text{morphisms} & &
 \end{array}$$

Thm (Vicary)

1) Every finite-dimensional C^* -algebra has a canonical structure of a special commutative ~~symmetric~~ dagger Frobenius algebra.

2) The category of commutative special dagger Frobenius and an appropriate notion of morphism is equivalent to the category of finite sets and functions.