

# Grothendieck 2-topoi are elementary 2-topoi

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# Subobject classifiers

## Definition.

Let  $\mathcal{C}$  be a 1-category with finite limits. A **subobject classifier** in  $\mathcal{C}$  is a **monomorphism**  $\tau: 1 \hookrightarrow \Omega$  in  $\mathcal{C}$  that is **universal** in the following sense:

$$\begin{array}{ccc} A & \longrightarrow & 1 \\ \downarrow \text{vi} & \lrcorner & \downarrow \tau \\ X & \overset{\exists! \chi_i}{\dashrightarrow} & \Omega \end{array}$$

For every  $X \in \mathcal{C}$  the function

$$\mathcal{G}_{\tau, X}: \mathcal{C}(X, \Omega) \rightarrow \text{Sub}(X)$$

given by **pulling back  $\tau$  is a bijection.**

# Elementary toposes

An **elementary topos** is a category  $\mathcal{C}$  with finite limits that has a subobject classifier and is cartesian closed.

**An elementary topos has an internal logic!**

## Example.

The archetypal elementary topos is *Set*.

$T: 1 \hookrightarrow \{T, F\}$  classifies subsets via the characteristic functions.

$$\begin{array}{ccc} A & \longrightarrow & 1 \\ \downarrow v_{i_A} & \lrcorner & \downarrow T \\ X & \overset{\exists! \chi_A}{\dashrightarrow} & \{T, F\} \end{array}$$

$$\chi_A : X \longrightarrow \{T, F\}$$

$$x \mapsto \begin{cases} T & \text{if } x \in A \\ F & \text{if } x \notin A \end{cases}$$

# 2-classifiers

## Definition (Weber, slightly changed).

Let  $\mathcal{K}$  be a 2-category with finite 2-limits. A **2-classifier** is a morphism  $\tau: 1 \rightarrow \Omega$  in  $\mathcal{K}$  such that for every  $X \in \mathcal{K}$  the functor

$$G_{\tau, X}: \mathcal{K}(X, \Omega) \rightarrow \mathcal{D}OpFib_{\mathcal{K}}(X)$$

given by **taking comma objects from  $\tau: 1 \rightarrow \Omega$**  is an **equivalence of categories**.

# Discrete opfibrations

We upgrade the subobjects to **discrete opfibrations**, that (essentially) **have as fibres arbitrary sets** (thus of one dimension higher).

## Example.

The archetypal elementary 2-topos is *CAT*.

**1:  $1 \rightarrow \mathcal{S}et$**  classifies all discrete opfibrations with small fibres, via the **category of elements (Grothendieck construction)**.

$$\begin{array}{ccc} \mathcal{E} & \longrightarrow & 1 \\ \downarrow \text{disc opfib} & \swarrow \text{comma} & \downarrow 1 \\ \mathcal{B} & \overset{\exists \chi_p}{\dashrightarrow} & \mathcal{S}et \\ \text{small fibres} & \text{taking fibres} & \end{array}$$

equivalence of categories

$$[\mathcal{B}, \mathcal{S}et] \simeq \mathcal{D}OpFib(\mathcal{B})$$

# Dense generators

## Definition.

A 2-functor  $J: \mathcal{Y} \rightarrow \mathcal{K}$  is **dense** if the restricted Yoneda embedding

$$\begin{aligned} \tilde{J}: \mathcal{K} &\longrightarrow [\mathcal{Y}^{\text{op}}, \mathcal{CAT}] \\ F &\mapsto \mathcal{K}(J(-), F) \end{aligned} \quad \text{is fully faithful.}$$

If  $J$  is fully faithful, this is equivalent to ask that **each object  $F \in \mathcal{K}$  is a weighted 2-colimit of objects of  $\mathcal{Y}$  which is preserved by  $\tilde{J}$ .**

## Example.

**Representables are a dense generator for 2-presheaves.** Every 2-presheaf is a weighted 2-colimit of representables.

# Reduction to dense generators

**We reduce the study of 2-classifiers to dense generators.**

Let  $J: \mathcal{Y} \rightarrow \mathcal{K}$  be a fully faithful dense 2-functor. Consider a discrete opfibration  $\tau: \Omega_{\bullet} \rightarrow \Omega$  in  $\mathcal{K}$ .

## Proposition (M.).

If  $\forall Y \in \mathcal{Y}$  the functor  $\mathcal{G}_{\tau, Y}: \mathcal{K}(Y, \Omega) \rightarrow \mathcal{D}OpFib(Y)$  is **faithful**, then  $\forall F \in \mathcal{K}$  the functor  $\mathcal{G}_{\tau, F}: \mathcal{K}(F, \Omega) \rightarrow \mathcal{D}OpFib(F)$  is **faithful**.

## Theorem (M.).

If  $\forall Y \in \mathcal{Y}$  the functor  $\mathcal{G}_{\tau, Y}: \mathcal{K}(Y, \Omega) \rightarrow \mathcal{D}OpFib(Y)$  is **fully faithful**, then  $\forall F \in \mathcal{K}$  the functor  $\mathcal{G}_{\tau, F}: \mathcal{K}(F, \Omega) \rightarrow \mathcal{D}OpFib(F)$  is **full** (and faithful).

# Reduction to dense generators

## Theorem (M.).

Assume that  $\forall Y \in \mathcal{Y}$  the functor  $\mathcal{G}_{\tau, Y}: \mathcal{K}(Y, \Omega) \rightarrow \mathcal{D}OpFib(Y)$  is an **equivalence of categories**. If an operation of normalization is possible, then  $\forall F \in \mathcal{X}$  the functor  $\mathcal{G}_{\tau, F}: \mathcal{K}(F, \Omega) \rightarrow \mathcal{D}OpFib(F)$  is **essentially surjective on objects** (and fully faithful).

$$\begin{array}{ccccc} H^{C, X} & \longrightarrow & G & & \Omega \bullet \\ \mathcal{G}_{\varphi}(\Lambda_{(C, X)}) \downarrow & \lrcorner & \downarrow \varphi & & \downarrow \tau \\ D(C, X) & \xrightarrow{\Lambda_{(C, X)}} & F & \xrightarrow{\chi^{\varphi}} & \Omega \\ & \searrow & \text{---} & \text{---} & \\ & & \mathcal{G}_{\tau}^{-1}(\mathcal{G}_{\varphi}(\Lambda_{(C, X)})) & & \end{array}$$



# Reduction applied to $\mathcal{CAT}$

## Example (Reduction of the category of elements).

The singleton category  $1$  is a dense generator in  $\mathcal{CAT}$ . So we can just look at the discrete opfibrations over  $1$ .

$$\mathcal{G}_{\tau,1}: \mathcal{CAT}(1, \mathit{Set}) \rightarrow \mathit{Set}$$

sends a functor  $1 \rightarrow \mathit{Set}$  to the set it picks, so it is an equivalence of categories. By the theorems of reduction, the construction of the **category of elements is fully faithful and classifies all discrete opfibrations with small fibres.**

Moreover, following the proof, we obtain a classifying morphism for a  $\varphi$  by collecting all its fibres, since the pullback of  $\varphi$  along  $B: 1 \rightarrow \mathcal{B}$  gives precisely the fibre over  $B$ .

# A 2-classifier in prestacks

Let  $\mathcal{C}$  be a category and consider  $\mathcal{K} = [\mathcal{C}^{\text{op}}, \mathcal{CAT}]$ . **Representables form a dense generator**, so we can just look at

$$\mathcal{G}_{\tau, y(\mathcal{C})}: [\mathcal{C}^{\text{op}}, \mathcal{CAT}](y(\mathcal{C}), \Omega) \rightarrow \mathcal{DOPFib}(y(\mathcal{C}))$$

and ask this to be an equivalence of categories for every  $C \in \mathcal{C}$ .

$\mathcal{C} \xrightarrow{\Omega} \mathcal{DOPFib}(y(\mathcal{C}))$  would only give a pseudofunctor  $\mathcal{C}^{\text{op}} \rightarrow \mathcal{CAT}$ .

Thanks to the **indexed Grothendieck construction (joint work with Elena Caviglia)**, we can take the strict replacement

$$\mathcal{C} \xrightarrow{\tilde{\Omega}} [(\mathcal{C}/\mathcal{C})^{\text{op}}, \mathit{Set}]$$

This is in line with Hofmann-Streicher.

# A 2-classifier in prestacks

## Theorem (M.).

Consider  $\tau: 1 \rightarrow \Omega$  such that  $\tau_C$  picks  $\text{id}_{y(C)} \in \mathcal{D}OpFib(y(C))$ . Then

$$\mathcal{G}_{\tau, y(C)}: \text{Ps}[C^{\text{op}}, \mathcal{CAT}](y(C), \Omega) \rightarrow \mathcal{D}OpFib(y(C)) = \Omega(C)$$

is **isomorphic to Yoneda's lemma map and is thus an equivalence.**

The corresponding  $\tilde{\tau}: 1 \rightarrow \tilde{\Omega}$  such that  $\tilde{\tau}_C$  picks  $\Delta 1$  is a **2-classifier in  $[C^{\text{op}}, \mathcal{CAT}]$ .**

Idea of the normalization process:

$$\begin{array}{ccc} H^{C,X} & \longrightarrow & G \\ \psi^{C,X} \downarrow & \lrcorner & \downarrow \varphi \\ y(C) & \xrightarrow{X} & F \end{array}$$

change the fibre  $(\psi_D^{C,X})_{D \rightarrow C}$   
into the fibre  $(\varphi_D)_{F(f)(X)}$   
(fibres of  $\varphi$  are global).

## Definition (Idea).

A **stack** is a **bicategorical sheaf**. Matching families are only required to satisfy the **descent compatibility up to isomorphism**.

A **descent datum** for  $F: \mathcal{C}^{\text{op}} \rightarrow \mathcal{CAT}$  w.r.t. a sieve  $S$  on  $C \in \mathcal{C}$  is

$$(D \xrightarrow{f} C) \in S \mapsto x_f \in F(D)$$

and for every  $D'' \xrightarrow{h} D' \xrightarrow{g} D \xrightarrow{f} C$  with  $f \in S$  an isomorphism

$$\varphi^{f,g}: g^* x_f \xrightarrow{\cong} x_{f \circ g}$$

$$\begin{array}{ccc} h^*(g^* x_f) & \xrightarrow{h^* \varphi^{f,g}} & h^*(x_{f \circ g}) \\ \parallel & & \downarrow \varphi^{f \circ g, h} \\ (g \circ h)^*(x_f) & \xrightarrow{\varphi^{f, g \circ h}} & x_{f \circ g \circ h} \end{array}$$

# A 2-classifier in stacks

We want to restrict the 2-classifier

$$\begin{aligned} \tilde{\Omega} : \mathcal{C}^{\text{op}} &\longrightarrow \mathcal{CAT} \\ \mathcal{C} &\mapsto [(C/C)^{\text{op}}, \text{Set}] \end{aligned}$$

in  $[\mathcal{C}^{\text{op}}, \mathcal{CAT}]$  to a **2-classifier  $\Omega_J$  in functorial stacks  $\mathcal{St}(\mathcal{C}, \mathcal{J})$** .

$\Omega_J$  needs to be **tight enough to be a stack** but at the same time **loose enough to host the classification**.

$$\begin{array}{ccccc}
 H & & i(H) & \xrightarrow{\quad} & \tilde{\Omega} \\
 \downarrow \varphi & & \downarrow i(\varphi) & \lrcorner & \downarrow \tilde{\tau} \\
 y(\mathcal{C}) & & i(y(\mathcal{C})) & \xrightarrow{\chi^{i(\varphi)}} & \tilde{\Omega} \\
 & & \downarrow i(\tau_J) & \lrcorner & \\
 & & i(\Omega_J) & \xrightarrow{j} & \\
 \exists i(\chi_J^\varphi) & \dashrightarrow & & & 
 \end{array}
 \quad \mathcal{St}(\mathcal{C}, \mathcal{J}) \xrightarrow{i} [\mathcal{C}^{\text{op}}, \mathcal{CAT}]$$

# A 2-classifier in stacks

## Theorem (M.).

$$\begin{aligned}\Omega_J : \mathcal{C}^{\text{op}} &\longrightarrow \mathcal{CAT} \\ \mathcal{C} &\mapsto \mathit{Sh}(\mathcal{C}/\mathcal{C})\end{aligned}$$

is a **functorial stack**.

For discrete opfibrations  $\varphi : \mathcal{H} \rightarrow y(\mathcal{C})$  in  $\mathit{St}(\mathcal{C}, \mathcal{J})$  over representables, the **characteristic morphism  $\chi^{l(\varphi)}$  factorizes through  $i(\Omega_J)$** .

And the normalization process is done in prestacks.

Then  $\tau_J : \Omega_{J,\bullet} \rightarrow \Omega_J$  is a **2-classifier in functorial stacks  $\mathit{St}(\mathcal{C}, \mathcal{J})$** , classifying all discrete opfibrations with small fibres.

# References



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