

A brief journey from Bicategories to Cartesian linear bicategories

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Bicategories

Jean Bénabou - 1967

Cartesian Bicategories

Carboni and Walters - 1987

Carboni, Kelly, Walters, and Wood - 2008

Linear Bicategories

Cockett, Koslowski, and Seely - 2000

Cartesian Linear Bicategories

- Bicategories are a weak version of 2-categories, where the horizontal composition is associative up to coherent isomorphism, and similarly for identities.
- The definition of cartesian bicategory axiomatizes important examples of bicategories such as **Rel**, the category of sets and relations.
- Linear bicategories are a generalization of bicategories which introduce a second composition, and the two compositions are linked by linear distributions as suggested by linear logic.
- Our goal is to extend the definition of cartesian bicategory to the linear setting.
- **Rel** is not just a bicategory, it is a linear bicategory. So it should be our first example of a cartesian linear bicategory.

A *bicategory* \mathcal{B} consists of

- A collection of objects (0-cells): A, B, C, \dots
- For each pair of objects (A, B) , a (hom-) category $\mathcal{B}(A, B)$ with

- Objects (1-cells): $A \xrightarrow{p} B$,
- Arrows (2-cells): $A \begin{array}{c} \xrightarrow{p} \\ \Downarrow \alpha \\ \xrightarrow{q} \end{array} B$

- Vertical composition \circ : $A \begin{array}{c} \xrightarrow{p} \\ \Downarrow \alpha \\ \xrightarrow{q} \\ \Downarrow \beta \\ \xrightarrow{r} \end{array} B \mapsto A \begin{array}{c} \xrightarrow{p} \\ \Downarrow \beta \circ \alpha \\ \xrightarrow{r} \end{array} B$

- For each object A , an identity functor

$$I_A : 1 \rightarrow \mathcal{B}(A, A)$$

- A horizontal composition (bi)functor ;

$$;_{A,B,C} : \mathcal{B}(A, B) \times \mathcal{B}(B, C) \rightarrow \mathcal{B}(A, C)$$

$$A \xrightarrow{f} B \xrightarrow{g} C \mapsto f;g : A \longrightarrow C$$

$$\begin{array}{c}
 \begin{array}{ccccc}
 A & \xrightarrow{f} & B & \xrightarrow{g} & C \\
 \downarrow \alpha & & \downarrow \beta & & \\
 A & \xrightarrow{f'} & B & \xrightarrow{g'} & C
 \end{array} \\
 \Downarrow \alpha; \beta
 \end{array}
 \mapsto
 \begin{array}{c}
 \begin{array}{ccc}
 A & \xrightarrow{f;g} & C \\
 \downarrow \alpha; \beta & & \\
 A & \xrightarrow{f';g'} & C
 \end{array}
 \end{array}$$

where it is associative up to coherent isomorphism, and similarly has two unitors.

- 0-cells are sets: A, B, C, \dots
- Each (hom-) category $\mathcal{B}(A, B)$ is $\mathcal{P}(A \times B)$ which is a poset under inclusion with
 - 1-cells are relations $R : A \rightarrow B$
 - 2-cells $A \begin{array}{c} \xrightarrow{R} \\ \Downarrow \alpha \\ \xrightarrow{R'} \end{array} B$ are $\alpha : R \subseteq R'$
- Composition (bi)functor is an ordinary composition between relations.

$$;_{A,B,C} : \mathcal{B}(A, B) \times \mathcal{B}(B, C) \rightarrow \mathcal{B}(A, C)$$

$$A \xrightarrow{R} B \xrightarrow{S} C \mapsto R;S := \{(a, c) \mid \exists b \in B \ (a, b) \in R \text{ and } (b, c) \in S\}$$

$$A \begin{array}{c} \xrightarrow{R} \\ \Downarrow \alpha \\ \xrightarrow{S} \end{array} B \begin{array}{c} \xrightarrow{R'} \\ \Downarrow \beta \\ \xrightarrow{S'} \end{array} C \mapsto A \begin{array}{c} \xrightarrow{R;S} \\ \Downarrow \alpha;\beta \\ \xrightarrow{R';S'} \end{array} C \quad (R;S \subseteq R';S')$$

- Identity $I_A : 1 \rightarrow \mathcal{B}(A, A)$

Example (Q-Rel where Q is a quantale)

First we recall definition of a quantale:

Definition (Quantale)

A quantale is a complete lattice Q which carries a monoid structure with neutral element e with an associative binary operation $*$: $Q \times Q \rightarrow Q$, satisfying a distributive property

$$a * \bigvee_{i \in I} b_i = \bigvee_{i \in I} (a * b_i), \quad \bigvee_{i \in I} a_i * b = \bigvee_{i \in I} (a_i * b)$$

For example: The quantale generated by a monoid $\mathcal{M} = (M, e, \cdot)$ is $\mathcal{PM} = (\mathcal{PM}, \subseteq, \{e\}, \cdot)$, where \mathcal{PM} is the power set of M .

Example (Q -Rel where Q is a quantale)

- 0-cells are sets A, B, C, \dots
- Since we can order Q -relations pointwise by using the quantale ordering

$$R \subseteq R' \iff \forall a, b \ R(a, b) \leq R'(a, b)$$

each (hom-) category $\mathcal{B}(A, B)$ is a poset under inclusion with

- 1-cells are Q -relations $R : A \rightarrow B$ ($R : A \times B \rightarrow Q$)
 - 2-cells are $\alpha : R \subseteq R'$
- Composition (bi)functor is

$$;_{A,B,C} : \mathcal{B}(A, B) \times \mathcal{B}(B, C) \rightarrow \mathcal{B}(A, C)$$

$$(R, S) \mapsto R;S(a, c) := \bigvee_{b \in B} R(a, b) * S(b, c)$$

$$\begin{array}{c} A \begin{array}{c} \xrightarrow{R} \\ \Downarrow \alpha \\ \xrightarrow{S} \end{array} B \begin{array}{c} \xrightarrow{R'} \\ \Downarrow \beta \\ \xrightarrow{S'} \end{array} C \mapsto A \begin{array}{c} \xrightarrow{R;S} \\ \Downarrow \alpha;\beta \\ \xrightarrow{R';S'} \end{array} C \quad (R;S \subseteq R';S') \end{array}$$

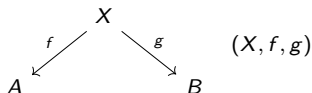
- Identity $I_A : 1 \rightarrow \mathcal{B}(A, A)$ where $I : A \rightarrow A$ is

$$I(a, a') = \begin{cases} e & \text{if } a = a' \\ \perp & \text{if } a \neq a' \end{cases}$$

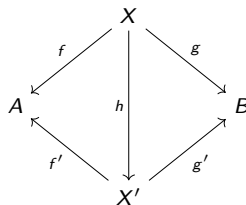
Example $\text{Span}(\mathcal{C})$ where \mathcal{C} is any category with pullbacks

- 0-cells are objects of \mathcal{C}
- Categories $\text{Span}(A, B)$ with

- 1-cells are $(f, g) : A \rightarrow B$



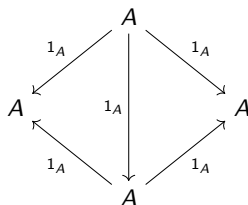
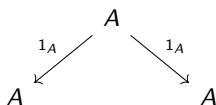
- 2-cells are $h : X \rightarrow X'$



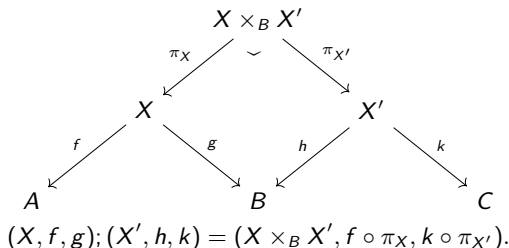
such that $f' \circ h = f$ and $g' \circ h = g$. Composition in $\text{Span}(A, B)$ is inherited from \mathcal{C} .

Example $\text{Span}(\mathcal{C})$

- Identity $I_A : 1 \rightarrow \text{Span}(A, A)$



- Composition (bi)functor is defined by using pullback as follows



Horizontal composition of 2-cells is defined by the universal property of pullbacks.

Homomorphism between bicategories

Let \mathcal{B} and \mathcal{B}' are two bicategories. A homomorphism $T : \mathcal{B} \rightarrow \mathcal{B}'$ consists of

- A function T mapping objects of \mathcal{B} into objects of \mathcal{B}'
- For each pair (A, B) of objects, a functor $T_{A,B} : \mathcal{B}(A, B) \rightarrow \mathcal{B}'(TA, TB)$
- For each triple (A, B, C) of objects, a natural transformation

$$\begin{array}{ccc}
 \mathcal{B}(A, B) \times \mathcal{B}(B, C) & \xrightarrow{i_{A,B,C}} & \mathcal{B}(A, C) \\
 T_{A,B} \times T_{B,C} \downarrow & \nearrow m_{A,B,C} & \downarrow T_{A,C} \\
 \mathcal{B}(TA, TB) \times \mathcal{B}'(TB, TC) & \xrightarrow{i_{TA,TB,TC}} & \mathcal{B}'(TA, TC)
 \end{array}$$

- For each object A , a natural transformation

$$u : l'_{TA} \Rightarrow T_{A,A} \circ l_A$$

satisfying coherence axioms.

Remark

In this section, bicategory \mathcal{B} is a locally posetal bicategory.

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Definition

A tensor product in \mathcal{B} is a homomorphism of bicategories

$$\otimes : \mathcal{B} \times \mathcal{B} \rightarrow \mathcal{B}$$

together with an object I , called the identity object, and natural isomorphisms

$$\rho : X \rightarrow X \otimes I$$

$$\gamma : X \otimes Y \rightarrow Y \otimes X$$

$$\alpha : (X \otimes Y) \otimes Z \rightarrow X \otimes (Y \otimes Z)$$

satisfying a number of coherence conditions. That is, we get a symmetric monoidal structure on \mathcal{B} .

A *cartesian* structure on a bicategory \mathcal{B} consists of

- A tensor product in \mathcal{B}
- On every object X in \mathcal{B} , a cocommutative comonoid structure. That is, arrows

$$\Delta_X : X \rightarrow X \otimes X, \quad t_X : X \rightarrow I.$$

These data must satisfy the following axioms

(U) Each arrow $r : X \rightarrow Y$ is a lax comonoid homomorphism. That is,

$$r; \Delta_Y \leq \Delta_X; (r \otimes r) \quad \text{and} \quad r; t_Y \leq t_X$$

$$\begin{array}{ccc} X & \xrightarrow{\Delta_X} & X \otimes X \\ \downarrow r & \leq & \downarrow r \otimes r \\ Y & \xrightarrow{\Delta_Y} & Y \otimes Y \end{array}$$

$$\begin{array}{ccc} X & & \\ \downarrow r & \searrow t_X & \\ Y & \xrightarrow{t_Y} & I \end{array}$$

(M) Comultiplication Δ_X and counit t_X have right adjoints Δ_X^*, t_X^* .

Definition

An arrow $r : X \rightarrow Y$ in a bicategory \mathcal{B} is called a *map* if it has a right adjoint r^* . Denote $\mathbf{Map}(\mathcal{B})$ the subcategory of \mathcal{B} determined by these maps.

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Theorem (Carboni and Walters)

Let \mathcal{B} be a locally posetal bicategory. If \mathcal{B} has a Cartesian structure, then

- i $\mathbf{Map}(\mathcal{B})$ has finite products,
- ii Each hom-category $\mathcal{B}(X, Y)$ has finite products which is denoted by \wedge .
- iii For any 1-cells r and s we have the following formula in \mathcal{B} :

$$r \otimes s = (p^*; r; p) \wedge (q^*; s; q) \quad (p \text{ and } q \text{ are appropriate projections})$$

Conversely, if $\mathbf{Map}(\mathcal{B})$ satisfies (i) and (ii) and the formula in (iii) defines a functorial tensor product on \mathcal{B} , then \mathcal{B} has a cartesian structure.

A *linear bicategory* \mathcal{B} consists of

- A collection of objects (0-cells): A, B, C, \dots
- For each pair of objects (A, B) , a (hom-) category $\mathcal{B}(A, B)$
- Two composition (bi)functors

$$;_{A,B,C}, \bullet_{A,B,C} : \mathcal{B}(A, B) \times \mathcal{B}(B, C) \rightarrow \mathcal{B}(A, C)$$

and units $\top_A, \perp_A \in \mathcal{B}(A, A)$, each coherently associate and unital. That is, we have two bicategory structures $(; , \top)$ and (\bullet , \perp) .

- Natural transformations δ_L and δ_R , called *linear distributivities*

$$\begin{array}{ccc}
 \mathcal{B}(A, B) \times \mathcal{B}(B, C) \times \mathcal{B}(C, D) & \xrightarrow{Id \times \bullet} & \mathcal{B}(A, B) \times \mathcal{B}(B, D) \\
 \downarrow ; \times Id & \swarrow \delta_L & \downarrow ; \\
 \mathcal{B}(A, C) \times \mathcal{B}(C, D) & \xrightarrow{\bullet} & \mathcal{B}(A, D)
 \end{array}$$

and,

$$\begin{array}{ccc}
 \mathcal{B}(A, B) \times \mathcal{B}(B, C) \times \mathcal{B}(C, D) & \xrightarrow{\bullet \times Id} & \mathcal{B}(A, B) \times \mathcal{B}(B, D) \\
 \downarrow Id \times ; & \nearrow \delta_R & \downarrow ; \\
 \mathcal{B}(A, C) \times \mathcal{B}(C, D) & \xrightarrow{\bullet} & \mathcal{B}(A, D)
 \end{array}$$

These must satisfy several coherence conditions.

Example (Rel)

- 0-cells are sets, A, B, C, \dots
- (hom-) category $\mathcal{B}(A, B)$ is $\mathcal{P}(A \times B)$ which is a poset under inclusion with
 - 1-cells are relations $R : A \nrightarrow B$

- 2-cells $A \begin{array}{c} \xrightarrow{R} \\ \Downarrow \alpha \\ \xrightarrow{R'} \end{array} B$ are $\alpha : R \subseteq R'$

- Two Composition (bi)functors

$$;_{A,B,C} : \mathcal{B}(A, B) \times \mathcal{B}(B, C) \rightarrow \mathcal{B}(A, C)$$

$$A \xrightarrow{R} B \xrightarrow{S} C \mapsto R; S := \{(a, c) \mid \exists b \in B \ (a, b) \in R \wedge (b, c) \in S\}$$

$$A \begin{array}{c} \xrightarrow{R} \\ \Downarrow \alpha \\ \xrightarrow{R'} \end{array} B \begin{array}{c} \xrightarrow{S} \\ \Downarrow \beta \\ \xrightarrow{S'} \end{array} C \mapsto A \begin{array}{c} \xrightarrow{R;S} \\ \Downarrow \alpha;\beta \\ \xrightarrow{R';S'} \end{array} C \quad (R; S \subseteq R'; S')$$

$$\bullet_{A,B,C} : \mathcal{B}(A, B) \times \mathcal{B}(B, C) \rightarrow \mathcal{B}(A, C)$$

$$A \xrightarrow{R} B \xrightarrow{S} C \mapsto R \bullet S := \{(a, c) \mid \forall b \in B \ (a, b) \in R \text{ or } (b, c) \in S\}$$

$$R \bullet S \subseteq R' \bullet S'$$

Let \mathcal{B} and \mathcal{B}' are two linear bicategories. A linear functor $F : \mathcal{B} \rightarrow \mathcal{B}'$ consists of

- A function F mapping objects of \mathcal{B} into objects of \mathcal{B}'
- For each pair (A, B) of objects, two functors $F_;, F_\bullet : \mathcal{B}(A, B) \rightarrow \mathcal{B}'(FA, FB)$
- 2-cells $m_\top : \top_{F_;(A)} \Longrightarrow F_{;(\top_A)}$ and $n_\perp : F_{\bullet;\perp_A}(A) \Longrightarrow \perp_{F_\bullet(A)}$
- Natural transformations, which with m_\top and n_\perp make $F_;$ monoidal (or "lax") with respect to $;$ and F_\bullet comonoidal (or "colax") with respect to \bullet

$$m_; : F_;(A);F_;(B) \Longrightarrow F_;(A;B)$$

and similarly for \bullet .

- Natural transformations (called "linear strengths")

$$\nu_;^R : F_;(A \bullet B) \Longrightarrow F_\bullet(A) \bullet F_;(B)$$

$$\nu_;^L : F_;(A \bullet B) \Longrightarrow F_;(A) \bullet F_\bullet(B)$$

and similarly for \bullet .

satisfying some coherence axioms.

The following can be seen as a lemma or definition of linear adjunction.

Lemma (Cockett, Koslowski, Seely)

Given 1-cells $A : X \rightarrow Y$, $B : Y \rightarrow X$ in a linear bicategory, the following are equivalent

- 1 A is left linear adjoint to B
- 2 For all 0-cells Z , the functor $(-);A : \text{Hom}(Z, X) \rightarrow \text{Hom}(Z, Y)$ is left adjoint to the functor $(-) \bullet B$
- 3 For all 0-cells Z , the functor $B;(-) : \text{Hom}(X, Z) \rightarrow \text{Hom}(Y, Z)$ is left adjoint to the functor $A \bullet (-)$

Remark

In this section, linear bicategory \mathcal{B} is a locally posetal linear bicategory.

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Definition

A tensor product and a cotensor product in \mathcal{B} are linear functors of linear bicategories

$$\otimes, \oplus : \mathcal{B} \times \mathcal{B} \rightarrow \mathcal{B}$$

together with two symmetric monoidal structures, denoted $(\mathcal{B}, \otimes, ;, \top)$ and $(\mathcal{B}, \oplus, \bullet, \perp)$, respectively. And, for each triple (X, Y, Z) of objects, the category \mathcal{B} must also be equipped with two distributive natural transformations:

$$\delta_L : X \otimes (Y \oplus Z) \longrightarrow (X \otimes Y) \oplus Z$$

$$\delta_R : (X \oplus Y) \otimes Z \longrightarrow X \oplus (Y \otimes Z)$$

satisfying a number of coherence conditions.

A *cartesian* structure on a locally posetal linear bicategory \mathcal{B} consists of

- A tensor \otimes product and a cotensor product \oplus in \mathcal{B} .
- On every object X in \mathcal{B} , a cocommutative comonoid structure, and a commutative monoid structure. That is, arrows

$$\Delta_X : X \rightarrow X \otimes X, \quad t_X : X \rightarrow \top,$$

$$\nabla_X : X \oplus X \rightarrow X, \quad \epsilon_X : \perp \rightarrow X.$$

satisfy the following axioms

(U₁) Each arrow $r : X \rightarrow Y$ is a colax comonoid homomorphism. That is

$$r; \Delta_Y \leq \Delta_X; (r \otimes r) \quad \text{and} \quad r; t_Y \leq t_X$$

$$\begin{array}{ccc} X & \xrightarrow{\Delta_X} & X \otimes X \\ \downarrow r & \leq & \downarrow r \otimes r \\ Y & \xrightarrow{\Delta_Y} & Y \otimes Y \end{array}$$

$$\begin{array}{ccc} X & & \\ \downarrow r & \searrow t_X & \\ Y & \xrightarrow{t_Y} & I \end{array}$$

(U₂) Each arrow $r : X \rightarrow Y$ is a lax monoid homomorphism. That is

$$(r \oplus r) \bullet \nabla_Y \leq \nabla_X \bullet r \quad \text{and} \quad \epsilon_X \bullet r \leq \epsilon_Y$$

$$\begin{array}{ccc} X & \xleftarrow{\nabla_X} & X \oplus X \\ \downarrow r & \geq & \downarrow r \oplus r \\ Y & \xleftarrow{\nabla_Y} & Y \oplus Y \end{array}$$

$$\begin{array}{ccc} X & & \\ \downarrow r & \nwarrow \epsilon_X & \\ Y & \xleftarrow{\epsilon_Y} & I \end{array}$$

(M) Comultiplication Δ_X and counit t_X have right linear adjoints.

Example (**Rel**)

- Compositions $R; S := \{(x, z) \mid \exists y \in Y \ (x, y) \in R \wedge (y, z) \in S\}$ and $R \bullet S := \{(x, z) \mid \forall y \in Y \ (x, y) \in R \vee (y, z) \in S\}$
- $(\mathbf{Rel}, \otimes, :, \top)$ is a symmetric monoidal category
 - Define $\otimes : \mathbf{Rel} \times \mathbf{Rel} \rightarrow \mathbf{Rel}$ on objects $X \otimes Y := X \times Y$
 - Define \otimes on morphisms by $R \otimes S : X \otimes X' \rightarrow Y \otimes Y'$ where $(x, x')R \otimes S(y, y')$ iff $(x, y) \in R \wedge (x', y') \in S$
- $(\mathbf{Rel}, \oplus, \bullet, \perp)$ is a symmetric monoidal category
 - Define \oplus on objects $X \oplus Y := X \times Y$
 - Define \oplus on morphisms by $R \oplus S : X \oplus X' \rightarrow Y \oplus Y'$ where $(x, x')R \oplus S(y, y')$ iff $(x, y) \in R \vee (x', y') \in S$

- Two distributive natural transformations:

$$\delta_L : X \otimes (Y \oplus Z) \longrightarrow (X \otimes Y) \oplus Z$$

$$\delta_R : (X \oplus Y) \otimes Z \longrightarrow X \oplus (Y \otimes Z)$$

- For every object X we have a cocommutative comonoid and a commutative monoid structure. That is, arrows $\Delta_X : X \rightarrow X \otimes X$ and $t_X : X \rightarrow \top$ and their duals $\nabla_X : X \oplus X \rightarrow X$ and $\epsilon_X : \perp \rightarrow X$

- Inequalities

$$\begin{array}{ll} r; \Delta_Y \leq \Delta_X; (r \otimes r) & \text{and} \quad r; t_Y \leq t_X \\ (r \oplus r) \bullet \nabla_Y \leq \nabla_X \bullet r & \text{and} \quad \epsilon_X \bullet r \leq \epsilon_Y \end{array}$$

- Comultiplication Δ_X and counit t_X have right linear adjoints by using the recent Lemma.

- Find and study more examples of cartesian linear bicategories
 - Bicategory of relations $\mathbf{Rel}(\mathcal{E})$ in a regular category \mathcal{E} ?
 - Ordered objects and ideals?
- Will have a theorem similar to Carboni and Walters's describing cartesian structure on $\mathbf{Map}(\mathcal{B})$?
- Can we define cartesian linear bicategories in general not just locally posetal?



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Thank you!